Docket No. AUS920010915US1

DYNAMIC SIZING LOGIC FOR DUMP LIST GENERATION

BACKGROUND OF THE INVENTION

1. Technical Field:

5 systems and, in particular, to performing a scan dump in the event of a system error. Still more particularly, the present invention provides a method, apparatus, and program for dynamically sizing and allocating memory for dump list generation.

10 2. Description of Related Art:

Some servers use diagnostic code to indicate failing components in the event of a system error. In many cases, however, these diagnostic codes are not sufficient to determine the nature of the failure. In these cases, selected chip data is saved in the event of a failure. Engineers may use this data to determine machine state and further diagnose the problem.

The process of saving this chip data is called a scan dump. The data, referred to as scan data, may contain various hardware dump information elements, such as scan ring, trace array, cache contents, and cache directory contents. During a complete dump the total dump set is dumped. However, depending on the error condition or user configuration, the system may decide to dump only a portion of the total dump set. This is referred to as an abbreviated dump.

When a dump is executed, the diagnostic code allocates memory for a dump list. The diagnostic code then builds the table and includes a list of hardware

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dump information elements to be collected based on the type of dump to be performed. To do this, the dump routine maintains a set of static arrays containing constants, each of which represents a scan ring or trace array, for example, to be dumped. The routine then iterates through the arrays to build the scan dump list. Another function is called for every entry in the dump list. This function collects hardware dump data and saves it into memory.

This process of executing a dump has several disadvantages. The process has multiple points of maintenance. If the static arrays are modified, for example a particular array is made bigger or smaller by adding or deleting constants, the iteration loops break because the information they have on how many entries they have to iterate through is inaccurate. In addition, the amount of memory allocated for the dump list is calculated using a predefined constant. The function used to allocate memory space for the list would have to be updated every time the number of possible dump list entries is changed.

Furthermore, the above process of executing a dump may result in memory waste. The function used to allocate memory space determines an amount of memory that would be needed for a complete dump. However, there may be several abbreviated dump modes that do not require nearly as much space. Therefore, if the diagnostic code is in one of the abbreviated dump modes, the process may allocate much more memory than is actually used.

Thus, it would be advantageous to provide an improved method, apparatus, and program for dynamically sizing and allocating memory for dump list generation.

Docket No. AUS920010915US1

SUMMARY OF THE INVENTION

The present invention provides an improved process for executing a dump. The iteration loops are made "smart" by allowing them to determine how big the arrays 5 are on the fly and adjust their behavior accordingly. The process uses a function to calculate the amount of memory to allocate for the dump list based on the dump mode and array sizes. Thus, if the static arrays are modified to add or delete constants or the diagnostic code is in an abbreviated dump mode, the amount of memory to be allocated will be calculated accurately.

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Docket No. AUS920010915US1

BRIEF DESCRIPTION OF THE DRAWINGS

The novel features believed characteristic of the invention are set forth in the appended claims. The invention itself, however, as well as a preferred mode of use, further objectives and advantages thereof, will best be understood by reference to the following detailed description of an illustrative embodiment when read in conjunction with the accompanying drawings, wherein:

Figure 1 depicts a block diagram of an illustrative embodiment of a data processing system with which the present invention may advantageously be utilized;

Figure 2 is a block diagram illustrating a hardware dump process in accordance with a preferred embodiment of the present invention; and

Figure 3 is a flowchart illustrating the execution of a dump in accordance with a preferred embodiment of the present invention.

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Docket No. AUS920010915US1

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENT

Referring now to the drawings and in particular to

Figure 1, there is depicted a block diagram of an
illustrative embodiment of a data processing system with
which the present invention may advantageously be
utilized. As shown, data processing system 100 includes
processor cards 111a-111n. Each of processor cards
111a-111n includes a processor and a cache memory. For
example, processor card 111a contains processor 112a and
cache memory 113a, and processor card 111n contains
processor 112n and cache memory 113n.

Processor cards 111a-111n are connected to main bus 115. Main bus 115 supports a system planar 120 that contains processor cards 111a-111n and memory cards 123. The system planar also contains data switch 121 and memory controller/cache 122. Memory controller/cache 122 supports memory cards 123 that includes local memory 116 having multiple dual in-line memory modules (DIMMs).

Data switch 121 connects to bus bridge 117 and bus

20 bridge 118 located within a native I/O (NIO) planar 124.

As shown, bus bridge 118 connects to peripheral
components interconnect (PCI) bridges 125 and 126 via
system bus 119. PCI bridge 125 connects to a variety of
I/O devices via PCI bus 128. As shown, hard disk 136 may

25 be connected to PCI bus 128 via small computer system
interface (SCSI) host adapter 130. A graphics adapter
131 may be directly or indirectly connected to PCI bus
128. PCI bridge 126 provides connections for external
data streams through network adapter 134 and adapter card

30 slots 135a-135n via PCI bus 127.

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Docket No. AUS920010915US1

An industry standard architecture (ISA) bus 129 connects to PCI bus 128 via ISA bridge 132. ISA bridge 132 provides interconnection capabilities through NIO controller 133 having serial connections Serial 1 and Serial 2. A floppy drive connection 137, keyboard connection 138, and mouse connection 139 are provided by NIO controller 133 to allow data processing system 100 to accept data input from a user via a corresponding input device. In addition, non-volatile RAM (NVRAM) 140 provides a non-volatile memory for preserving certain types of data from system disruptions or system failures, such as power supply problems. A system firmware 141 is also connected to ISA bus 129 for implementing the initial Basic Input/Output System (BIOS) functions. A service processor 144 connects to ISA bus 129 to provide functionality for system diagnostics or system servicing. The operating system (OS) is stored on hard disk 136, which may also provide storage for additional application software for execution by data processing system. NVRAM 140 is used to store system variables and error information for field replaceable unit (FRU) isolation. During system startup, the bootstrap program loads the operating system and initiates execution of the operating system. To load the operating system, the bootstrap program first locates an operating system kernel type from hard disk 136, loads the OS into memory, and jumps to an initial address provided by the operating

30 processing system. Once loaded and initialized, the operating system controls the execution of programs and

into random-access memory (RAM) within the data

system kernel. Typically, the operating system is loaded

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may provide services such as resource allocation, scheduling, input/output control, and data management.

The present invention may be executed in a variety of data processing systems utilizing a number of different hardware configurations and software such as bootstrap programs and operating systems. The data processing system 100 may be, for example, a stand-alone system or part of a network such as a local-area network (LAN) or a wide-area network (WAN).

Diagnostic code may execute on service processor 144 to indicate failing components in the event of a system The diagnostic code may be stored in system firmware 141. The diagnostic code may execute a dump process to save diagnostic data from chips in the system, such as processors 112a, 112n and memory controller/cache In accordance with a preferred embodiment of the present invention, the diagnostic code dynamically calculates the amount of memory to allocate for a dump list based on the dump mode and array sizes.

The data processing system illustrated in Figure 1 is an example of a data processing system on which the present invention may be implemented. Those of ordinary skill in the art will appreciate that the hardware depicted in Figure 1 may vary. For example, rather than the single memory controller/cache 122, the data processing system may include a memory controller/cache on each memory card 123. The depicted example is not meant to imply architectural limitations with respect to the present invention.

With reference now to Figure 2, a block diagram 30 illustrating a hardware dump process is shown in accordance with a preferred embodiment of the present

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Docket No. AUS920010915US1

invention. Diagnostic code **210** indicates failing components in the event of a system error. The diagnostic code determines a dump mode and array sizes for the dump mode based on static arrays **212** maintained in the diagnostic code.

One static array is maintained for each type of component to be scanned. For example, if the dump mode calls for the diagnostic code to receive dump data from three processors and three memory controllers, static arrays 212 include a static array for processors and a static array for memory controllers. Each static array includes a constant for each hardware dump information element to be collected.

Based on the dump mode, diagnostic code 210

15 determines how many components are to be scanned, which static arrays correspond to the components, and which hardware dump information elements are to be scanned.

Using this information, the diagnostic code determines the number and size of arrays to include in dump list 230 and, thus, calculates the amount of memory to allocate for the dump list in memory 220.

Thereafter, the diagnostic code builds dump list 230 that includes a list of hardware information elements to be collected, such as scan rings, trace arrays, cache contents, and cache directory contents. The diagnostic code iterates through the arrays to build the dump list. Another function is called for every entry in the dump list. This function collects the hardware dump data from chips, such as chips 202, 204, 206, 208, and saves it in memory 220.

Turning now to **Figure 3**, a flowchart is shown illustrating the execution of a dump in accordance with a

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preferred embodiment of the present invention. The process begins, determines the dump mode (step 302), and determines array sizes (step 304). Next, the process calculates the amount of memory for the dump list (306)

and builds the dump list (step **308**). Thereafter, the process saves the hardware dump information elements into memory (step **310**) and ends.

Thus, the present invention solves the disadvantages of the prior art by providing an improved process for executing a dump. The iteration loops are made "smart" by allowing them to determine how big the arrays are on the fly and adjust their behavior accordingly. The process uses a function to calculate the amount of memory to allocate for the dump list based on the dump mode and array sizes. Therefore, if the static arrays are modified to add or delete constants or the diagnostic code is in an abbreviated dump mode, the amount of memory to be allocated will be calculated accurately.

It is important to note that while the present invention has been described in the context of a fully functioning data processing system, those of ordinary skill in the art will appreciate that the processes of the present invention are capable of being distributed in the form of a computer readable medium of instructions and a variety of forms and that the present invention applies equally regardless of the particular type of signal bearing media actually used to carry out the distribution. Examples of computer readable media include recordable-type media, such as a floppy disk, a hard disk drive, a RAM, CD-ROMs, DVD-ROMs, and transmission-type media, such as digital and analog communications links, wired or wireless communications

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Docket No. AUS920010915US1

links using transmission forms, such as, for example, radio frequency and light wave transmissions. The computer readable media may take the form of coded formats that are decoded for actual use in a particular data processing system.

The description of the present invention has been presented for purposes of illustration and description, and is not intended to be exhaustive or limited to the invention in the form disclosed. Many modifications and variations will be apparent to those of ordinary skill in the art. The embodiment was chosen and described in order to best explain the principles of the invention, the practical application, and to enable others of ordinary skill in the art to understand the invention for various embodiments with various modifications as are suited to the particular use contemplated.